Data Storage Solutions for Decentralized Online Social Networks

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Research @ SANDS

	recommendation and decision support systems					entralized online social orking and collaboration		Applications
	privacy aware/preserved data aggregation, storage, sharing & analytics/data-mining				distributed key-value stores			
					data-center design	 	P2P/F2F storage systems	tributed) S
data/computation at 3 rd party/outsourced & data management systems							0	ystems
	social trust network models analysis			secure/privacy preserved computation primitives			codes for storage	Foundational

- Selective information dissemination using social links
 - GoDisco

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- # P2P storage

http://SandS.sce.ntu.edu.sg/d SN

- Not the same as a file-sharing system
 - Peer-to-Peer (P2P) storage systems leverage the combined storage capacity of a network of storage devices (peers) contributed typically by autonomous end-users as a common pool of storage space to store content reliably.

Design space

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 - Sophisticated functionalities: Concurrency, Version Control, ...

Realizing Reliability



Redundancy Type

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Replication



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Replication 貒 Erasure codes **※ B**₁ Retrieve any $\mathbf{0}_{1}$ \mathbf{O}_1 k' (\geq k) blocks **B**₂ 0 0 Data = ObjectB Encoding Decoding 0 Lost blocks Original k blocks **B**_n k blocks n encoded blocks

(stored in storage devices in a network)

Data

Reconstruct

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 - Security & privacy implications of data placement ...



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 - may lead to poor performance
 - access latency, repair cost, ...





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♦ ???

• Cons of TotalRecall, which placed at random:



Cloud assisted storage system

Source: Google tech talk on Wuala: <u>http://www.youtube.com/watch?v=3xKZ4KGkQY8</u>

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 - So very few hops needed, gives high through-put



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Control

• De/centralized, local/global knowledge

Replica Placement in P2P Storage: Complexity and Game Theoretic Analyses Rzadca et al, ICDCS 2010

- Replication model: A clique of replicas storing each other's data (reciprocity)
 - Explores both centralized and decentralized settings for clique formation
 - Challenge
 - Centralized matching right set of peers to optimize storage capacity utilization (proven NP-hard)
 - Decentralized matching uses an underlying gossip algorithm (T-man) to explore partners

(simulations with artificial data)



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- Maps naturally to the overlying social application
 - Anecdotal note: SafeBook used Friend-of-Friends for access control also

Place data at friends: That's it?

- Store at all friends (naïve/baseline)
 - Best one can do in terms of achieving highest possible availability
 - Very high overheads!
 - Storage
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Sharma et al, P2P 2011

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 - Criticality of friends
 - Which friends are indispensable?

Evaluation

- Data set
 - Italian instant messenger service
 - + Pros
 - Social+Temporal characterisitcs
 - "May" reasonably reflect the online/offline behavior
 - + Cons:
 - Not a p2p storage system trace
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Data set

- Italian instant messenger service
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3436 nodes

- 848 nodes in the largest component
 - Note that many nodes had "neighbors" in other servers, for whom we did not have info.
 - Between 1-18 neighbors
- Use two weeks of data
 - \circ One for "learning", one for evaluation
 - Time of day, day of week effects
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 \odot 50% nodes can get more than 90% availability

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If there are "enough" friends, (>10), ought to be okay! (assuming storage capacity is not an issue)

Bootstrapping pangs!

- New peers with few friends in the system, or no reputation of being highly available, will find it difficult to get started!
 - Game-theoretic study on reciprocity based P2P cliques
 - Analysis of ego-centric networks for F2F storage

Sharma et al, Comsnets 2012

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 - Superpeers help coordinating, finding storage partners, etc.

Take with a huge pinch of salt: artificial data to drive simulations, with too many parameters ...



Moving forward

	Light weight P2P OSN		Full-fledged (D)OSN
Bulk (static) data storage	dynamic/social data store High availability High consistency High rate of data updates Small volume of data	Security modules Encryption access control 	Social modules Analytics Search/Navigation Recommendation
P2P overlay with basic services: DHT lookup, peer-sampling, etc.			
Could be even nulti-)cloud based.			

